



Education

The Benefits of Solid State in Enterprise Storage Systems

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Agenda

- Why flash in the datacenter? Why now?
- Memory, cache and storage
- Application opportunities
- Flash in enterprise storage today
 - ◆ SSD storage tier
 - ◆ Storage controller-based cache
 - ◆ Flash in host systems
- What's next
- Conclusion

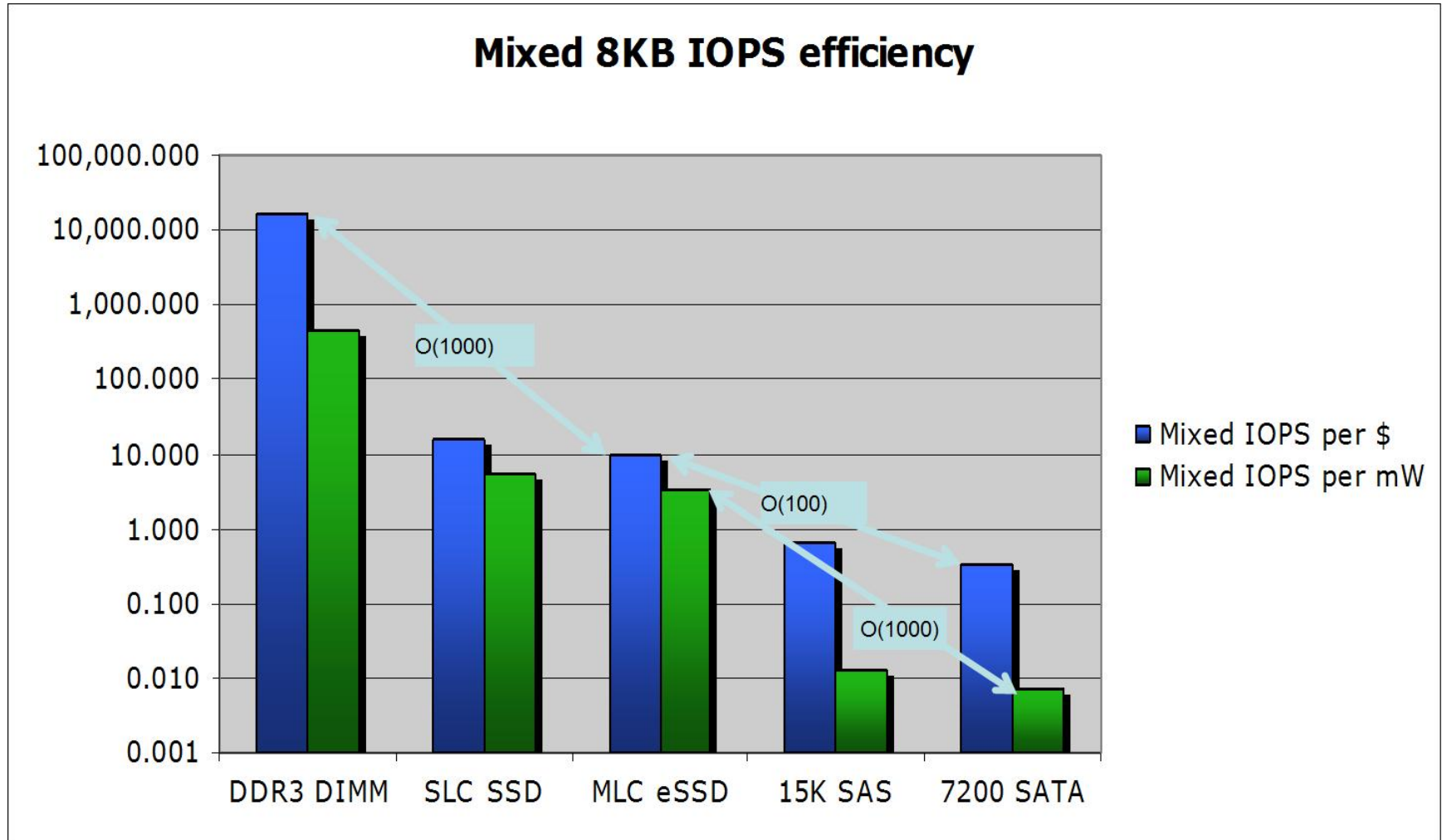
➤ Why flash?

- ◆ Capacity efficiency versus DRAM
 - › ~5x better \$ per GB
 - › ~40x better power per GB
- ◆ IOPS efficiency versus HDDs
 - › ~40x better \$ per IOPS
 - › ~600x better power per IOPS

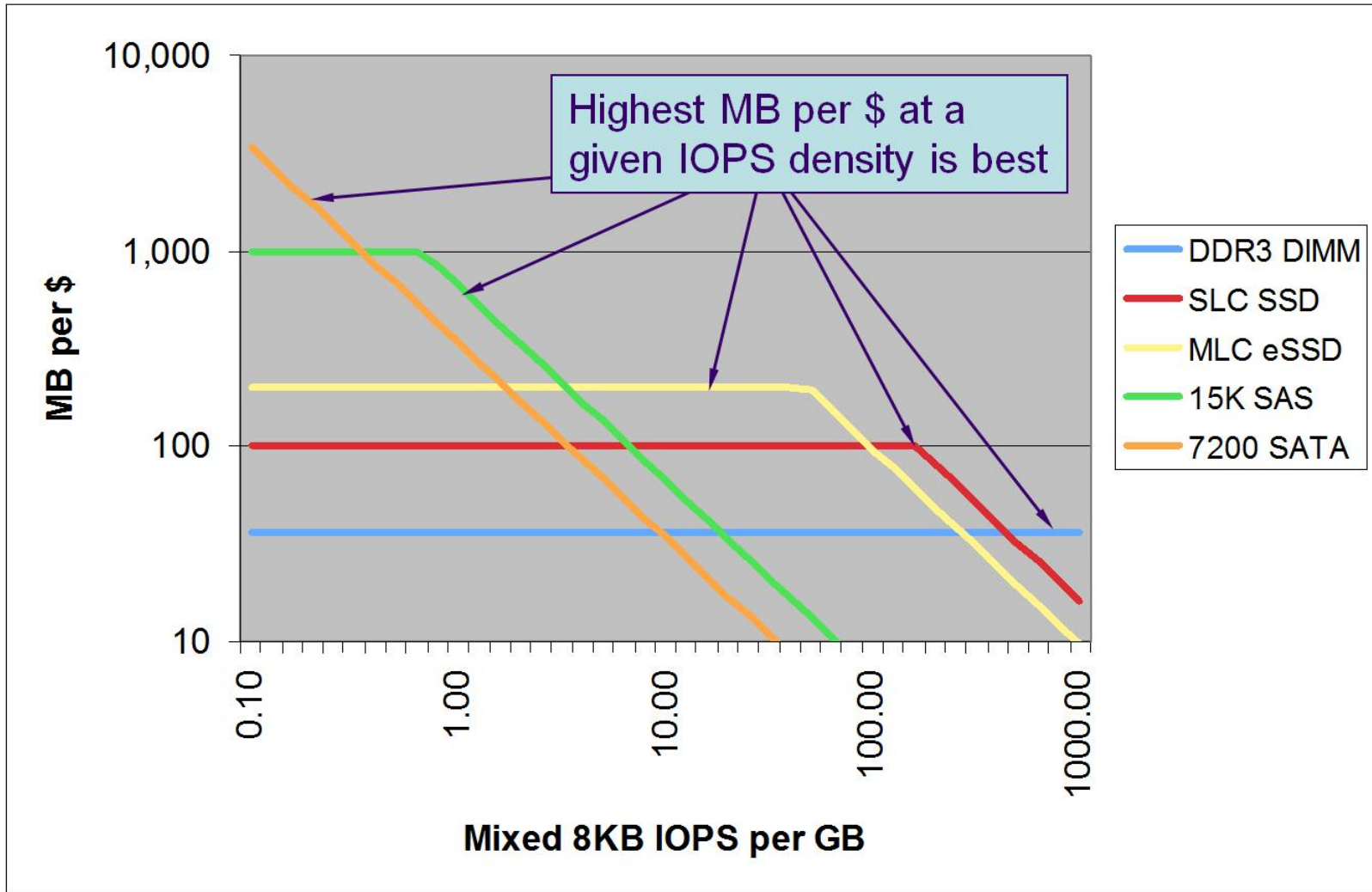
➤ Why now?

- ◆ Period of rapid density advancements led to HDD-like bit density at lower \$/GB than DRAM
- ◆ Innovations in SSD and tiering technology

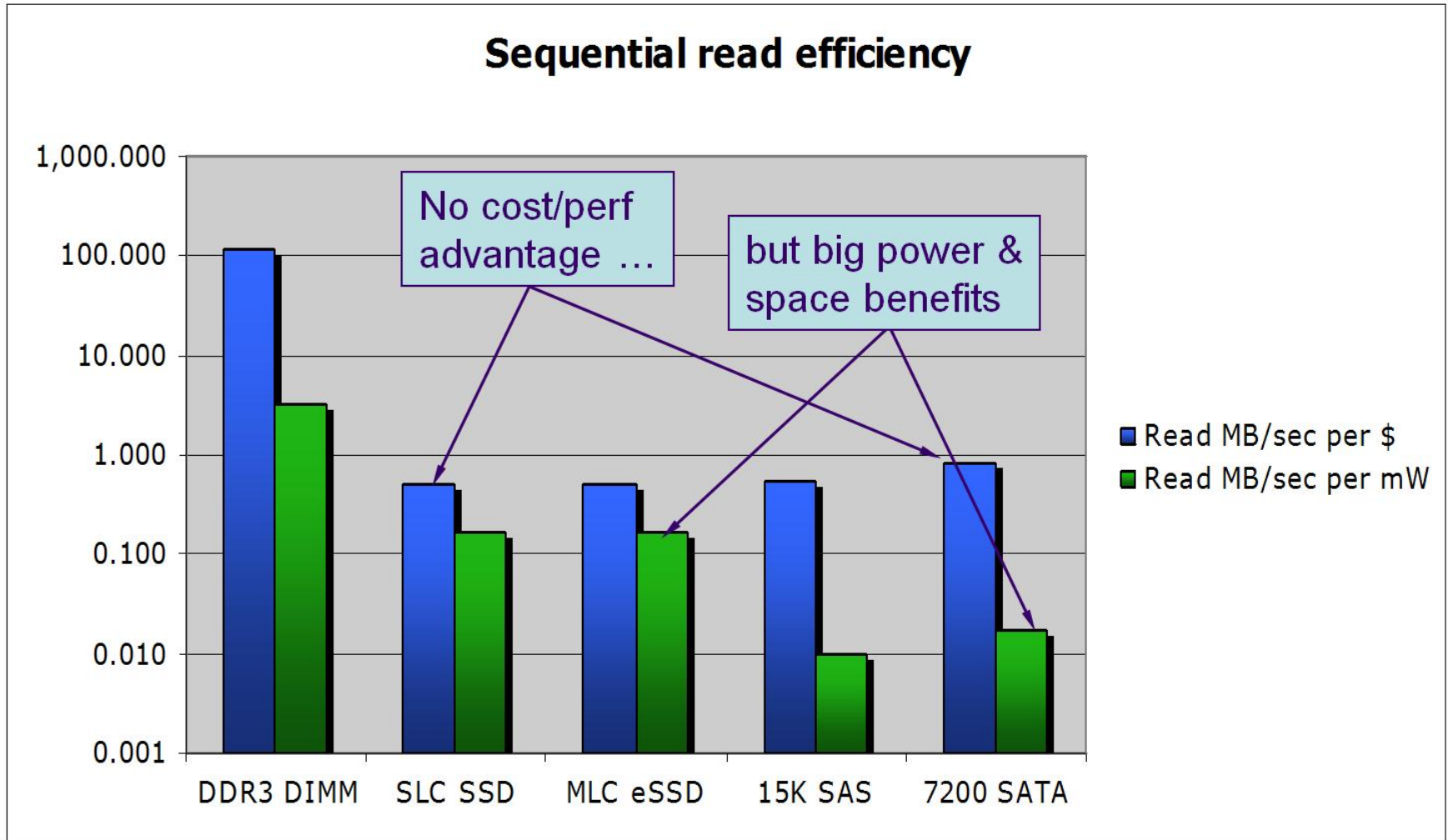
Why Flash? IOPS Efficiency



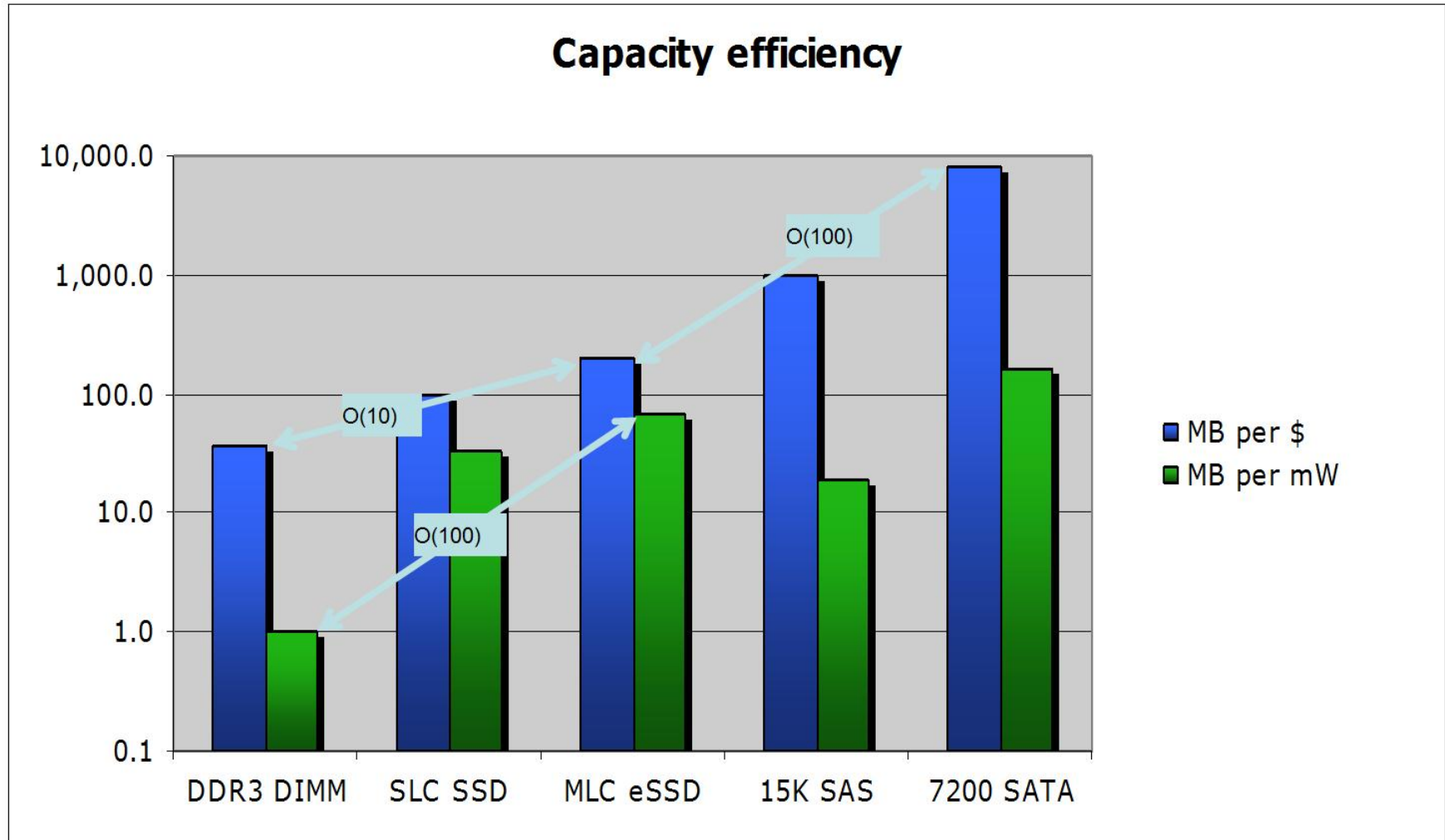
Why Flash? An IOPS Density View



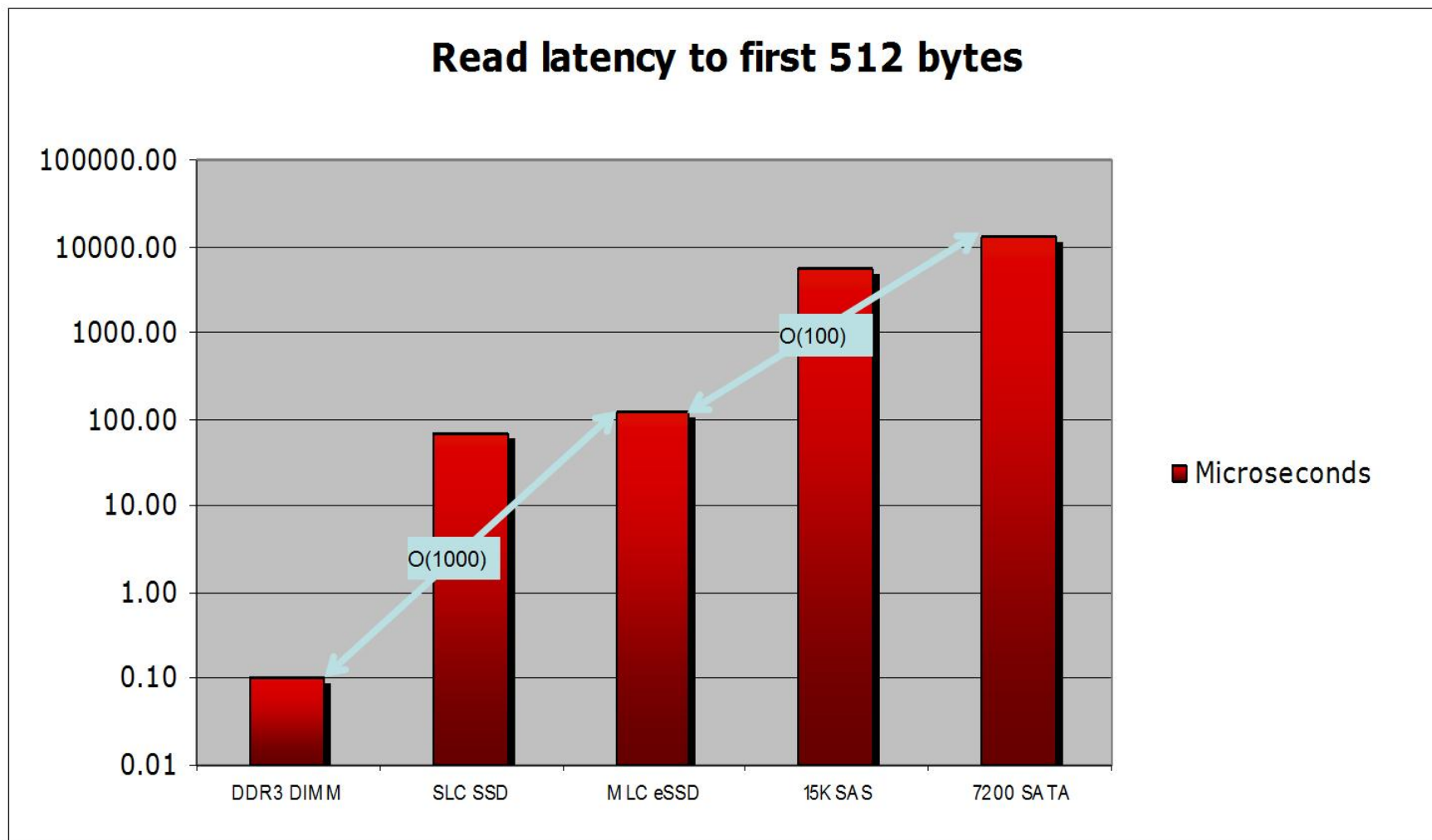
Why Flash? Read Throughput per Watt



Why Flash? Capacity Efficiency

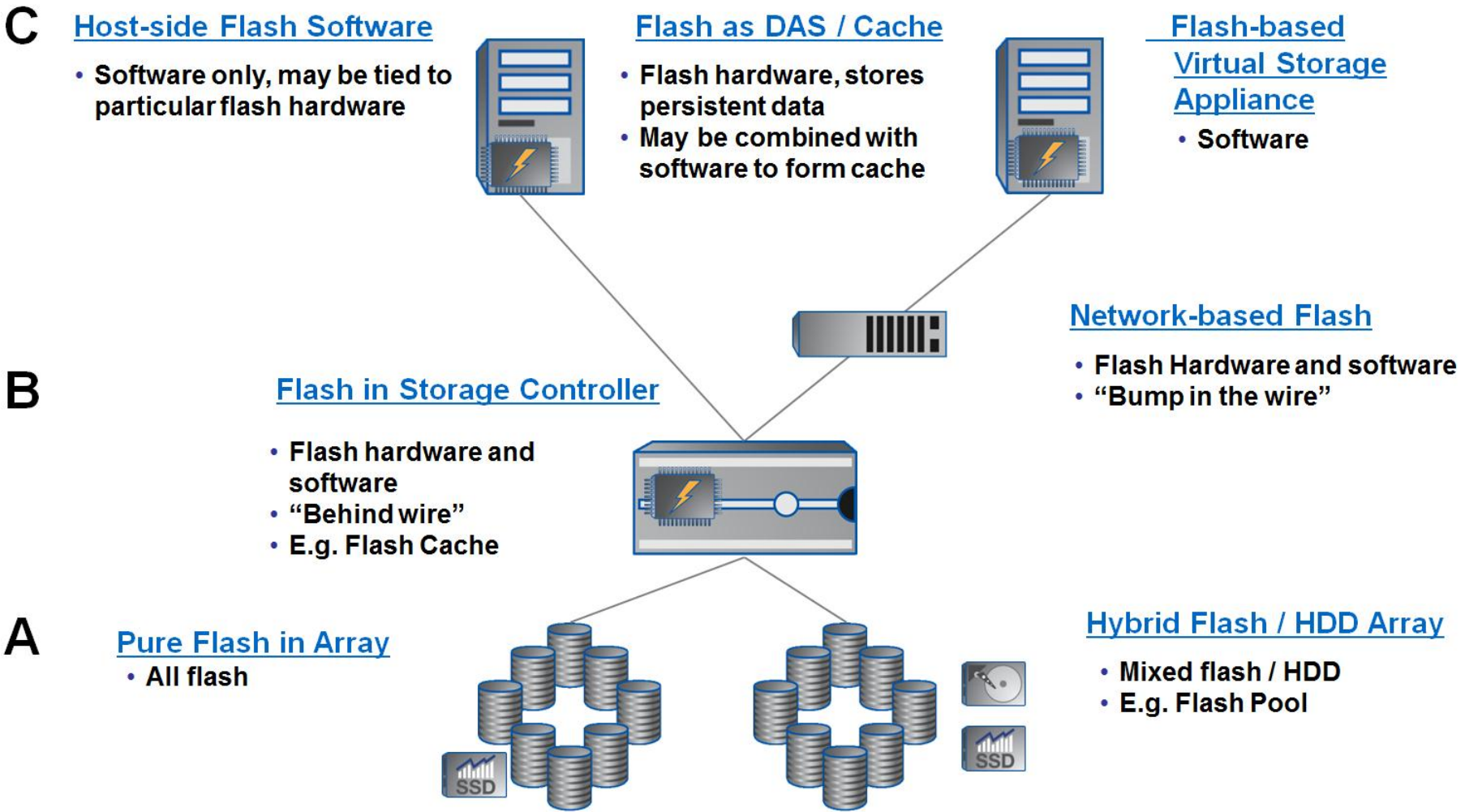


Why Flash? Read Latency



- Intense random reads, e.g. OLTP, metadata
- Sequential read after random write
 - ◆ Log-oriented writes convert this to random read after sequential write (e.g. FTL)
- Low read latency (~100x better than HDD)
 - ◆ Facilitates DRAM extension by allowing high read throughput with limited read concurrency
 - ◆ Paging datacenter apps can be practical again
 - ◆ Memory capacity to consolidate more servers with underutilized CPU
- Enabling memory-resident datasets, e.g.
 - ◆ OLTP
 - ◆ Data warehouses (*viz* TPC-H results)
 - ◆ Large metadata

End-to-End Flash Categories



(A) Solid State Disk Tier

➤ Advantages:

- ◆ Fast random I/O for small blocks
- ◆ Low read and write latency time
- ◆ Low power consumption
- ◆ Low noise
- ◆ Better mechanical reliability

➤ Disadvantages:

- ◆ Very high price, typically 10-30 X comparable FC drives
- ◆ Limited capacities
- ◆ Slow random write speeds, e.g. erase of blocks
- ◆ Slow sequential write throughput

➤ Database acceleration solution

- ◆ Entire database on SSD tier, or
- ◆ Hot random access files on SSD and rest of database on standard disk
 - › Indexes and temp space

➤ Large scale virtual machine environments

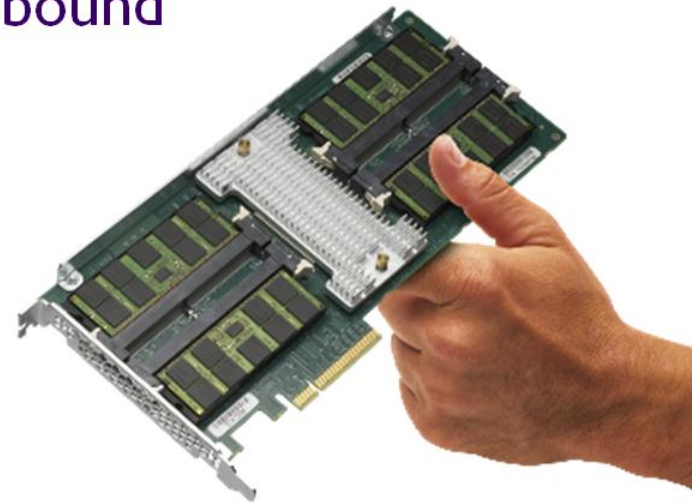
- ◆ Solves “boot storm” problem for large numbers of virtual machines
- ◆ Deduplication of VM data, e.g. virtual desktops
 - › Reduces capacity requirements, increasing IOPS density, potentially making SSD economical

Automated Tiering or Tier-less

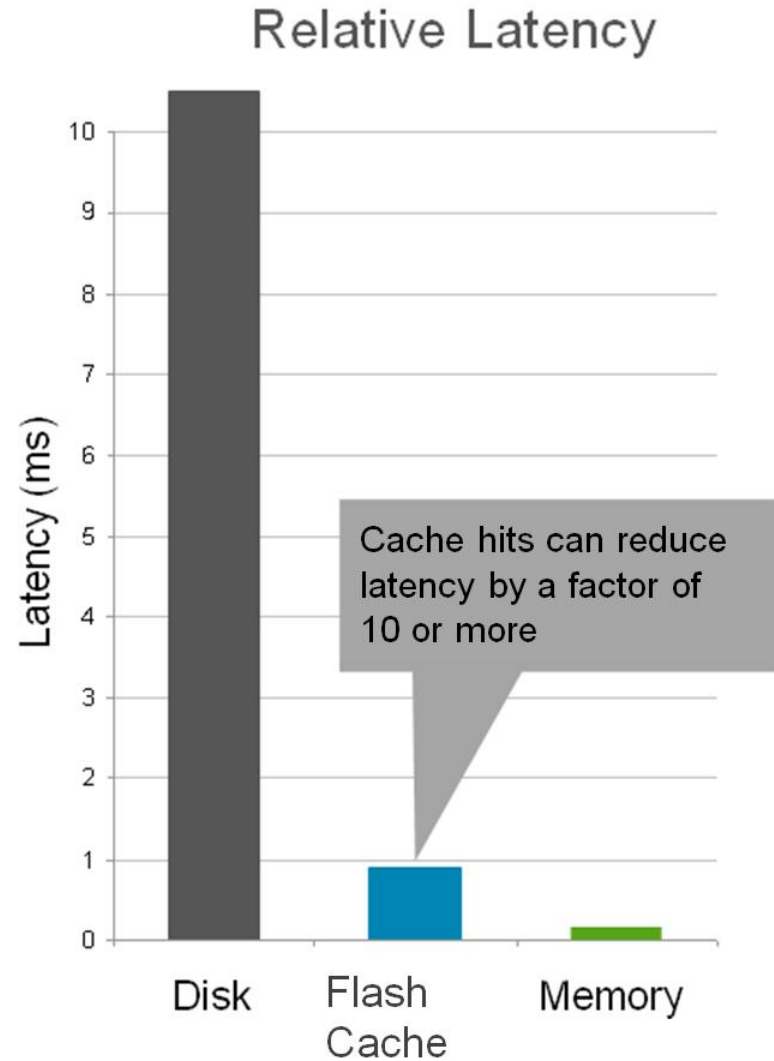
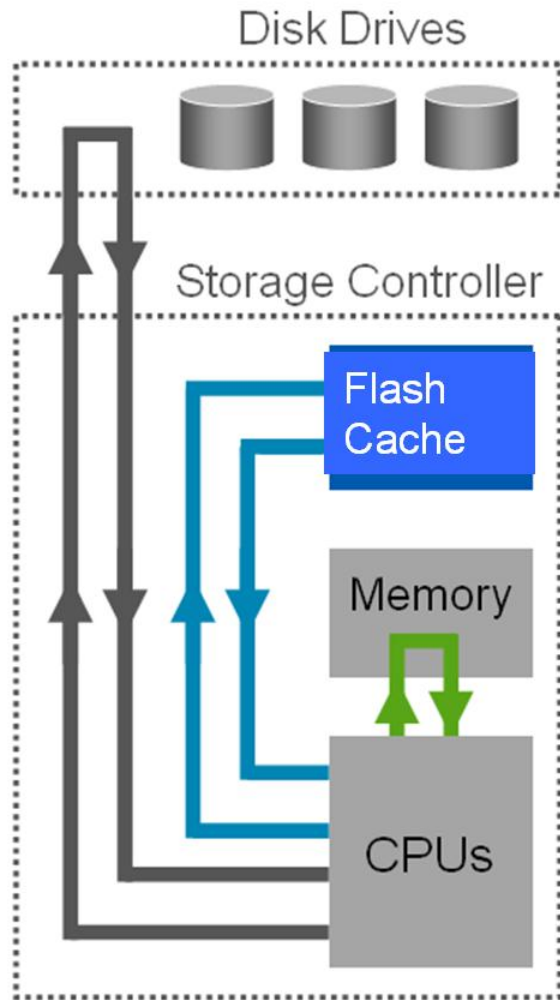
- Mixing SSD and HDD for a particular workload will probably be the most cost-efficient use of SSDs in over the next few years
- Area of intense innovation among enterprise storage vendors
- Issue is to automate data placement and movement
 - ◆ Automated tiering
 - ◆ Policy-based
 - ◆ No administrator overhead imposed
 - ◆ Some vendors refer to this as tier-less storage
- As SSD prices fall this will become increasingly important

(B) Controller-based Flash Cache

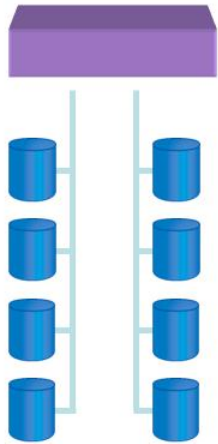
- Functions as an intelligent read cache for data and metadata
- Automatically places active data where access can be fast
- Provides more I/O throughput without adding high-performance disk drives to a disk-bound storage system
- Effective for file services, OLTP databases, messaging, and virtual infrastructure



Reduce Latency with Flash Cache

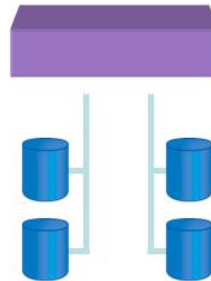


Use case: Scale Performance of Disk-bound Systems



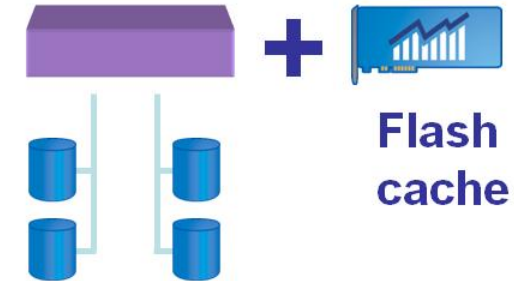
Add Spindles

- Use more disks to provide more IOPs
- May waste storage capacity
- Consumes more power and space



Starting Point: **Need More IOPs**

- Performance is disk-bound
- Have enough storage capacity
- Random read intensive workload

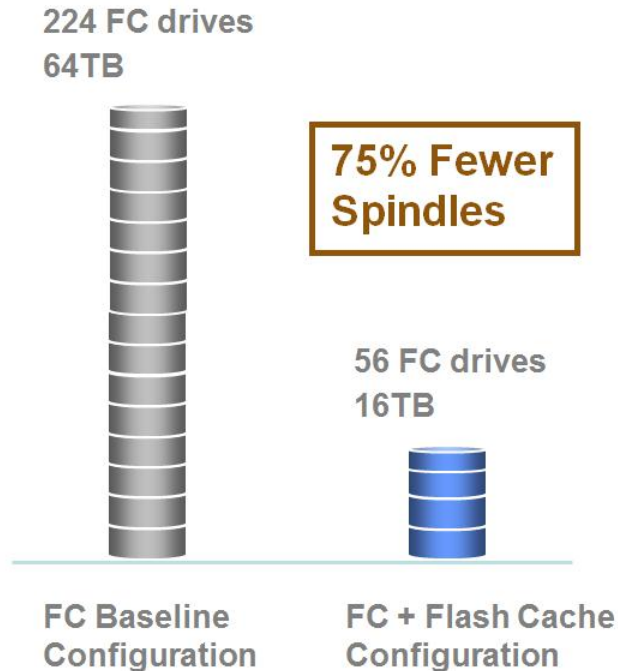


Add Flash Cache

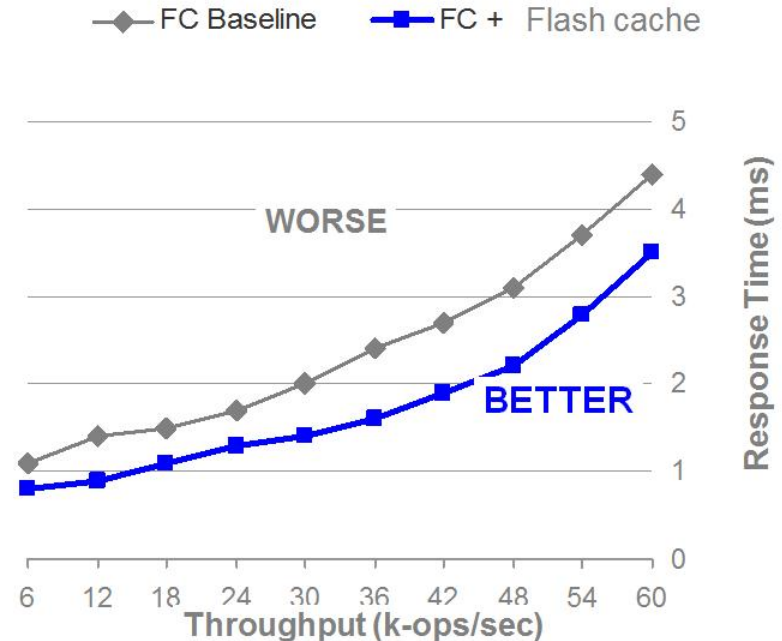
- Use cache to provide more IOPs
- Improves response times
- Uses storage efficiently
- Achieves cost savings for storage, power, and space

FC HDD plus Flash Cache Example

Benchmarked Configurations



SPECsfs2008 Performance



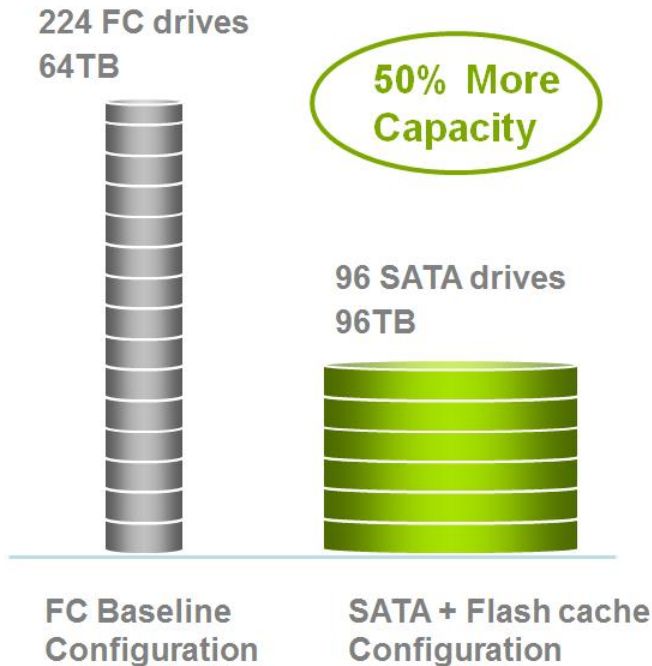
- Purchase price is **50% lower** for FC + Flash cache compared to Fibre Channel baseline
- FC + Flash cache yields **67% power savings** and **67% space savings**

For more information, visit <http://spec.org/sfs2008/results/sfs2008nfs.html>.

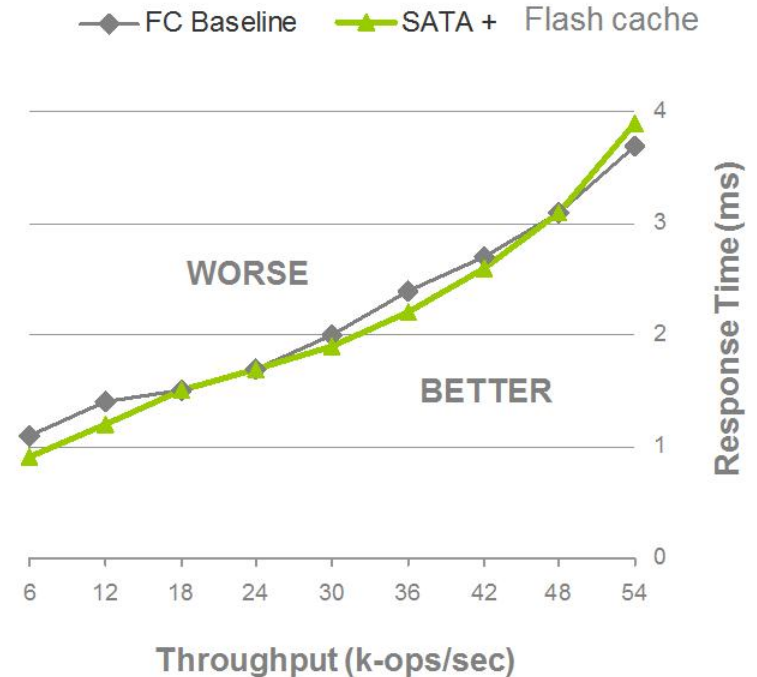
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SATA HDD plus Flash Cache Example

Benchmarked Configurations



SPECsfs2008 Performance



- Purchase price is **39% lower** for SATA + Flash cache compared to FC baseline
- SATA + Flash cache yields **66% power savings** and **59% space savings**

For more information, visit <http://spec.org/sfs2008/results/sfs2008nfs.html>.

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(C) Host-based Flash

- Flash card on PCI bus in host system
- Can support SCSI semantics or device driver model
- Acts as Tier 0 storage (IOPS tier) in front of networked storage (capacity tier)
 - ◆ Requiring no data movement (caching)
 - ◆ Requiring data movement (AST)
- Multiple implementations in development:
 - ◆ High performance DAS
 - ◆ Shared storage RAID subsystem in VM on host
 - ◆ Shared storage OS in VM on host
- Area of intense industry and standards activity

Cost Structure of Memory/Storage Technologies

Cost determined by

- cost per wafer
- # of dies/wafer
- memory area per die [sq. μm]
- memory density [bits per $4F^2$]
- patterning density [sq. μm per $4F^2$]

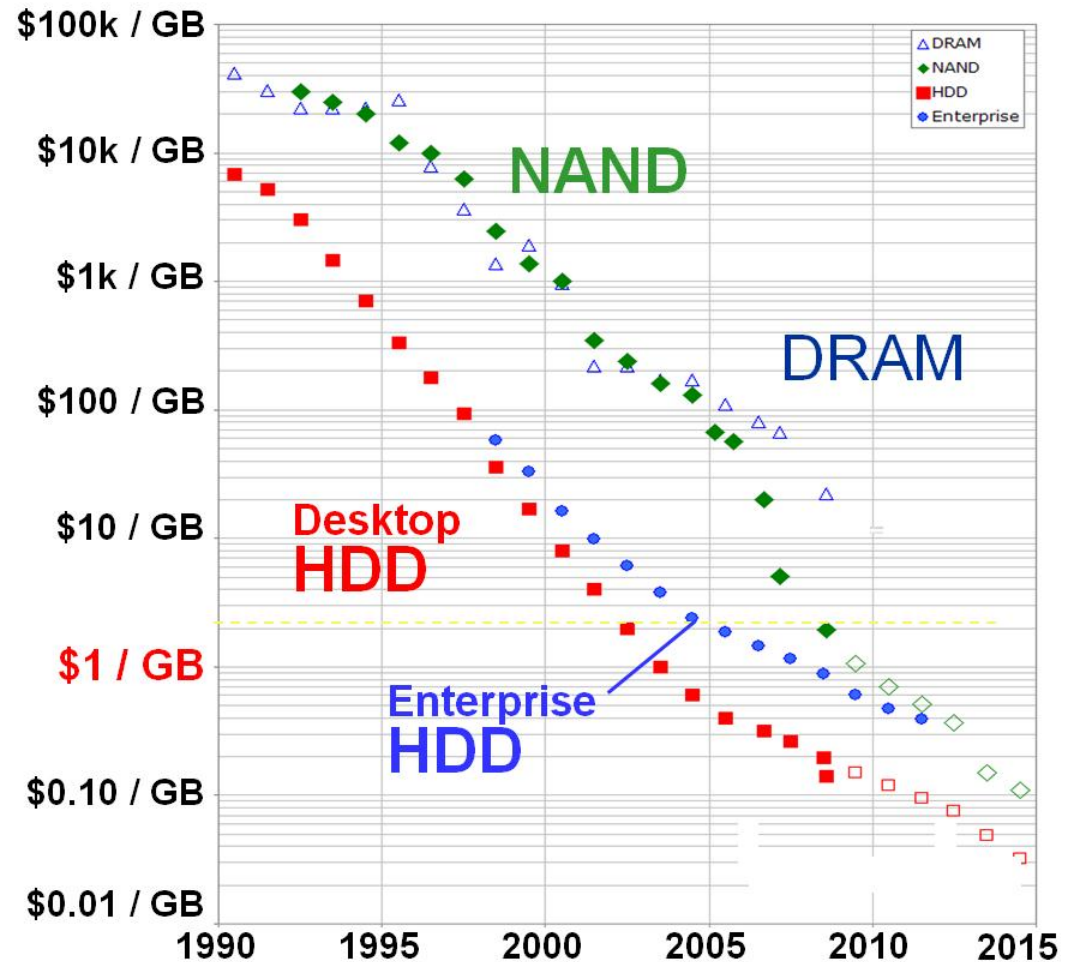
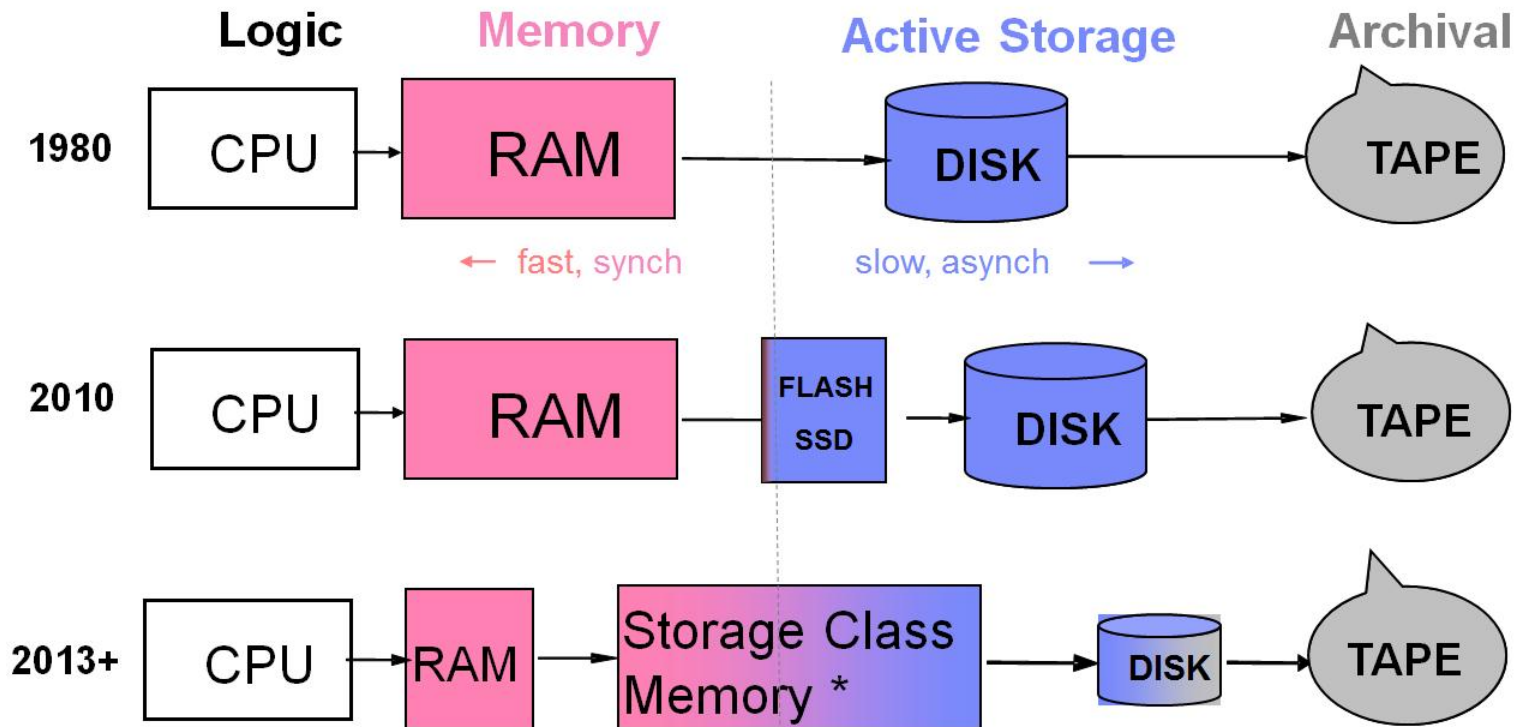


Chart courtesy of Dr. Chung Lam,
IBM Research updated version
of plot from 2008 *IBM Journal R&D* article

System Evolution



* e.g. Phase change memory
Memristor
Solid Electrolyte
Racetrack memory

- Over the next 5 years solid state technologies will have a profound impact on enterprise storage
- It's not just about replacing mechanical media with solid state media
- The architectural balance of memory, cache and persistent storage will change
- Today's solid state implementations in enterprise storage demonstrate these changes
- It's only the beginning...

The SNIA Education Committee would like to thank the following individuals for their contributions to this Tutorial.

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